

On Bayesian Approach to Grillage Optimization

Jonas Mockus¹, Rimantas Belevičius², Dmitrij Šešok³,
Juozas Kaunas⁴, Darius Mačiūnas⁵

¹ Department of Information Systems,
^{2,3,4,5} Department of Engineering Mechanics,
Faculty of Fundamental Sciences, Vilnius Gediminas Technical University
Saulėtekio av. 11, LT – 10223 Vilnius, Lithuania
e-mail: jonas.mockus@vgtu.lt¹, rimantas.belevicius@vgtu.lt², dmitrij.sesok@vgtu.lt³,
juozas.kaunas@gmail.com⁴, darius.maciunas@vgtu.lt⁵

crossref <http://dx.doi.org/10.5755/j01.itc.41.4.1670>

Abstract. In this paper, a new simplified version of the Bayesian Approach to coordinate global optimization (BAcoor) is compared with the well-known algorithms. BAcoor is a method of multi-dimensional optimization by applying a sequence of one-dimensional global optimizers starting from the best points obtained by previous one-dimensional optimization. The globality of one-dimension search is controlled by the only parameter. The new element is that observation points are defined by explicit formulas. In other similar methods this is performed by some numerical techniques that minimize the risk functions. The efficiency of suggested method is investigated and compared with other methods by solving a real-life civil engineering global optimization problem of pile placement schemes in grillage-type foundations. This problem is a good benchmark, because the minimal value of the objective function is known so the optimization error can be defined exactly.

Keywords: global optimization; finite element method; genetic algorithms; optimization of grillages.

1. Introduction

In this paper we suggest a new simplified algorithm BAcoor (Bayesian Approach to coordinate global optimization) performance of which is controlled by the only parameter, and compare it with other well-known algorithms. As the benchmark for the comparison, one particular engineering optimization problem, for which the global solution is known in advance, is chosen: optimization of pile placement schemes in grillage-type foundations. Grillage-type foundations (or simply “grillages”) are the most popular and effective scheme of foundations. Grillages consist of supporting piles and connecting beams, and transfer the dead loadings from the erection together with effective loading to the ground. If the minimum required number of piles is known (and in case of homogeneous piles it can be simply obtained dividing the total loading by the carrying capacity of a pile), the optimal placement scheme is characterized by an even distribution of reactive forces arising in piles. This can be achieved choosing appropriate positions of all piles. Thus, optimal design of pile placement schemes under grillages is a specific problem of civil engineering. However, it ideally suits for comparison of global optimization algorithms. First of all, the global

solution is known in advance, therefore we can estimate the achieved results. Secondly, our experience shows that the objective function for practical grillage optimization problems has complex surface and possesses many local minima points. Moreover, usually the objective function is very sensitive to the positions of piles: even a small alteration of one position leads to a significant change of the objective function. All this makes the optimal design of pile placement schemes under grillages a difficult global optimization problem.

Exhaustive technical details on the design of grillages can be found, e.g. in [14]. However, only a few papers deal with optimization of placement schemes. In [4] combination of the sizing and topology optimization is observed, however the piles are aggregated to special groups of piles. In [7] minimization of the differential settlements of piled rafts is analyzed, again, by a special way minimizing the number of design variables. We are trying to obtain the globally minimal price of pile foundations treating all piles as a separate design variables. In [1], [2] and [3] the idealizations of real grillages are introduced, which are taken in the present mathematical model as well. In the last paper the problem under discussion served as a benchmark for comparison of different

global optimization algorithms: modified random search meta-heuristics (genetic algorithms and simulated annealing), and local optimization combined with random search. Ten practical grillages of small-to-medium scale (17 to 55 design variables) were optimized; data on these grillages were obtained from the Dutch design bureaus. In this paper we optimize again the same grillages with the suggested algorithm BAcoor and compare obtained results.

2. Optimization algorithm

2.1. Outline

The traditional numerical analysis considers optimization algorithms that guarantee some accuracy for all functions to be optimized. To limit maximal errors, one needs computational efforts that often increase exponentially with the size of the problem. This is the main disadvantage of this approach. This is the main reason to justify the heuristic methods which developed using the expert knowledge. This includes meta-heuristics such as Genetic Algorithm (GA) [6], Simulated Annealing (SA) [9], and Simplex [12].

The alternative is the average case analysis. Here an average error is not limited but is made as small as possible. The average is taken over a set of functions to be optimized. In [5, 10], the average case analysis is called the Bayesian Approach (BA).

The Bayesian Approach (BA) is defined by fixing a prior distribution P on a set of functions $f(x)$ and by minimizing the Bayesian risk function $R(x)$ which defines approximately the expected deviation of $f(x)$ from the global minimum at a fixed point x . The distribution P is considered as a stochastic model of $f(x)$, $x \in R^m$, where $f(x)$ might be a deterministic or a stochastic function. The objective of BA is to provide as small average error as possible while keeping convergence conditions.

2.2. Simplified Bayesian Approach

In [8, 15, 16], the Wiener process was used as a stochastic model when $m = 1$. In [10, 11], the simplified models were designed and the simple expression of $R(x)$ was obtained by replacing the traditional Kolmogorov consistency conditions when $m > 1$. However, the simplified stochastic models [10, 11] can be preferable in one-dimensional cases, too, since they offer some computational advantages by providing explicit solutions and the Wiener model is just an approximation of actual functions $f(x)$.

The aim of BA is to reduce the expected deviation. In addition, BA has some good asymptotic properties, too. It is shown in [10] that

$$d^*/d_a = \left(\frac{f_a - f^* + \varepsilon}{\varepsilon} \right)^{\frac{1}{2}}, \quad n \rightarrow \infty. \quad (1)$$

Here d^* is a density of points x_i around the global optimum. d_a is an average density of x_i in the feasible area. f_a is an average value of $f(x)$ in this area. f^* is an average value of $f(x)$ around the global minimum. ε is the correction parameter, and n is a number of observations (function $f(x_i)$ evaluations). That means that BA provides convergence to the global minimum for any continuous $f(x)$ and a greater density of observations x_i around the global optimum, if n is large.

2.3. One-dimensional Optimization

If the number of variables $m = 1$, then

$$R(x) = \min_{1 \leq i \leq n} z_i - \min_{1 \leq i \leq n-1} \frac{(x - x_i)^2}{z_i - c_n}. \quad (2)$$

The minimum of the risk function $R(x)$ is obtained at the point

$$x_{n+1} = \arg \max_{a \leq x \leq b} \min_{1 \leq i \leq n-1} \frac{(x - x_i)^2}{z_i - c_n} \quad (3)$$

$$z_i = f(x_i), c_n = \min_{1 \leq i \leq n} z_i - \varepsilon.$$

Here the maximum is reached at points $x = x_i^*$, $i = 1, \dots, n-1$ satisfying conditions

$$\frac{(x - x_i)^2}{z_i - c_n} = \frac{(x - x_{i+1})^2}{z_{i+1} - c_n} \quad (4)$$

$$x_i \leq x \leq x_{i+1}, \quad x_1 = a, \quad x_n = b, \quad i = 1, 2, \dots, n-1. \quad (5)$$

This way condition (3) is reduced to

$$x_{n+1} = \arg \min_{1 \leq i \leq n-1} R(x_i^*) \quad (6)$$

here

$$R(x_i^*) = \min_{1 \leq i \leq n-1} z_i - \frac{(x_i^* - x_i)^2}{z_i - c_n}. \quad (7)$$

Fig.1 illustrates the simple case when $n = 5$ and, $x_6 = \arg \min_i R(x_i^*)$, $i = 1, 2, 3, 4$.

The points satisfying (3) are defined by conditions

$$x_i^+ = \frac{v_2 X_1 - v_1 X_2 + (X_2 - X_1) \sqrt{v_1 v_2}}{v_2 - v_1}, \quad (8)$$

$$x_i^- = \frac{v_2 X_1 - v_1 X_2 - (X_2 - X_1) \sqrt{v_1 v_2}}{v_2 - v_1}, \quad (9)$$

where $v_1 = z_i - c_n$, $v_2 = z_{i+1} - c_n$, $X_1 = x_i$, $X_2 = x_{i+1}$, $i = 1, \dots, n-1$.

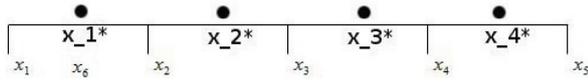


Figure 1. Illustrates the points of observations x_i , $i = 1, 2, 3, 4, 5$ and the solutions x_i^* , $i = 1, 2, 3, 4$, $x_1 = a$, $x_5 = b$, where a and b define the interval of feasible x .

The points satisfying both (3) and (4) are determined as follows

$$x_i^* = x_i^+, \text{ if } x_i \leq x_i^+ \leq x_{i+1} \quad (10.1)$$

$$x_i^* = x_i^-, \text{ if } x_i \leq x_i^- \leq x_{i+1} \quad (10.2)$$

According to expression (1), large ε defines global, nearly uniform search. Using small ε , most of the observations are performed around global minima. The global convergence is provided at any $\varepsilon > 0$.

2.4. Coordinate Optimization (BAcoor)

The one-dimensional model is applied for coordinate optimization when $x = (x^j, j = 1, \dots, m)$. We start from some initial point $x = x(0)$ and apply one-dimensional search by this sequential procedure:

The first cycle of BAcoor is defined by these expressions:

$$x^1(1) = \arg \min_{a_1 \leq x^1 \leq b_1} f(x^1, x^2(0), \dots, x^m(0)) \quad (11)$$

$$x^2(1) = \arg \min_{a_2 \leq x^2 \leq b_2} f(x^1(1), x^2, \dots, x^m(0)) \quad (12)$$

where $a_j \leq x^j \leq b_j$ are the feasible intervals. If some improvement is achieved, the second cycle starts in a similar way. The search stops if no improvement occurs.

The coordinate optimization converges to global minimum along each coordinate and is a convenient tool of visualization by providing a set of one-dimensional projections. However, to guarantee the global convergence, BAcoor should be repeated many times from different initial points $x(0)$.

3. Idealizations and mathematical formulation of benchmark problem

The initial data for the grillage optimization problem are as follows:

- The geometrical scheme of connecting beams;
- Cross-section data of all beams (area, moments of inertia);
- Material data of all beams (material in one beam is treated as isotropic);
- Positions of immovable piles (if any);
- Maximum allowable reactive force at any pile (all piles are homogeneous);
- Minimum possible distance between adjacent piles;

- Vertical and two rotational stiffnesses (along the beam and normal to the beam) of pile;
- Loading data. Active forces can be applied in the form of concentrated loads and moments at any point on beam, or in the form of distributed trapezoidal loadings at any segment of beam.

The results of optimization are appropriate positions of piles under connecting beams at which the reactive forces in piles do not exceed the carrying capacity of piles. In an ideal grillage, reactive forces at all piles are identical. Practically this is hardly feasible, particularly when a designer introduces the so-called “immovable supports” that have to retain their positions and cannot change them during optimization process. Some technological constraints may also make the ideal scheme non-achievable, for example, the distance between adjacent piles cannot be too small due to the specific capacities of a pile driver. In the present work we do not consider the immovable supports and allow for a pile to take whatever position in the grillage, thus typically the piles are not placed at the joints of grillage. This fact confines the pile placement problem scope to a low-rise buildings without significant overturning moments due to horizontal thrust, e.g. due to earthquake loading or wind loads.

The objective function for minimization can be formulated in several alternative forms, e.g. the maximal vertical reactive force at a pile, the difference between the maximal and minimal reactive forces in the whole grillage, or the maximal difference between the reactive force and carrying capacity of a pile (for piles with different carrying capacities). We assume that the characteristics of all piles are equal and use the first form of objective function:

$$f^* = \min_{x \in D} f(x) \quad (13)$$

where $f(x)$ is a nonlinear objective function of continuous variables $f: \mathfrak{R}^m \rightarrow \mathfrak{R}$, m is the number of design parameters x defining positions of piles, $D \subset \mathfrak{R}^m$ is a feasible region of design parameters. No assumptions on unimodality are included into formulation of the problem. The maximal vertical reactive force at a pile is considered as the objective function:

$$f(x) = \max_{i=1, \dots, N_a} F_i(x) \quad (14)$$

where N_a is the number of piles, $F_i(x)$ is the reactive force at the i -th pile.

Since a supporting pile may reside only under connecting beams, there are evident restrictions on the positions of piles: during the optimization process the piles can move only along the connecting beams. Therefore, a two-dimensional beam structure of the grillage is “unfolded” to a one-dimensional construct, and the piles are allowed to range through this space freely. Unfortunately, in such a formulation, small

variation of the design variable may correspond to a finite variation of the position of the pile in the physical space, what leads to discontinuity of the problem. One possibility to overcome this is to use multilevel optimization where the upper level combinatorial problem assigns piles to beams, while the lower level continuous problems aim to position the piles in the assigned beams. Another possibility is to divide search space avoiding jumps of piles from one beam to another and perform searches in such separate spaces. However we do not know what piles and how many piles to assign to particular connecting beams because sometimes very different topologies of placement schemes lead to close values of objective function. Thus, the parameterization that would not introduce discontinuities and lead to simpler optimization problem is not applicable to the problem.

One design parameter corresponds to a position of one pile in the one-dimensional construct ($n = N_a$). The backward transformation restores the positions of piles in the original beam structure of the grillage. The constraints for the design parameters are as follows:

$$0 \leq x_i \leq L, \quad i = 1, \dots, N_a \quad (15)$$

where x_i is a design parameter defining the position of the i -th pile. L is the total length of all beams in the grillage. If the minimal possible distance δ between adjacent piles is specified, there are additional constraints

$$\|x_i - x_j\| \geq \delta, \quad i \neq j \quad (16)$$

where x_i are two-dimensional coordinates of piles and $\|x_i - x_j\|$ denotes the distance between piles. To cope with this constraint, a penalty is included in the objective function.

A finite element program is used as a “black-box” routine to the optimization program for solution of direct problem to find reactive forces in the grillage. In the direct problem that is solved via finite element analysis, the connecting beams in the grillage are idealized as the beam elements, while the piles are treated as supports, i.e. finite element mesh nodes with given elastic boundary conditions. Since time of optimization crucially depends on time of solution of the direct problem, fast problem-oriented original FORTRAN programs with a special mesh pre-processor have been developed and used.

The beam elements have 2 nodes with 6 degrees of freedom each (3 displacements along the coordinate axes and 3 rotations about these axes). The expressions of structural matrices of element can be found in many textbooks on finite elements.

The main statics equation is

$$[K]\{u\} = \{P\} \quad (17)$$

where $[K]$ is the stiffness matrix of the ensemble of elements, $\{u\}$ are the nodal displacements, and $\{P\}$ are the active forces.

The reactive forces at piles are available after obtaining the nodal displacements:

$$F_i = \sum_j [K_{ij}] u_j \quad (18)$$

Sensitivity analysis may be used if an optimization algorithm requires information about derivatives.

4. Numerical results and discussion

The pile placement schemes of 10 practical grillages possessing from 17 to 55 piles (design variables) were optimized in [1] using 7 different optimization algorithms. Data for these problems (Appendix 1) are obtained from several Dutch design bureaus (courtesy of Consultancy W. F. O. B.V., Paauw B.V. Aannemingsbedrijf, Aannemingsbedrijf V. Dijk, Bouwtectuur West Friesland, Stabo Bouw B.V., Aannemingsbedrijf A. Tuin Den Helder and others) which use the professional software package *MatrixFrame* (<http://www.matrix-software.com/uk/structuralengineering/matrixframe/index.html>) for structural engineering. It is intended for an analysis and design of steel and concrete erections. All those problems chosen for comparison of algorithms have one common trait: the proportion between the total loading and the allowable reaction is such that the engineering solution (i.e., when the actual reactive forces do not exceed the allowable reaction at any pile) requires achieving almost the ideal solution.

In each algorithm the total number of objective function evaluations was the same: 5000. Three stochastic global optimization algorithms were investigated: the Modified Random Search (MRS) [1], Simulated Annealing (SA) [9], Genetic Algorithm (GA) [6]. In addition, three local methods with random initial points were regarded including the Simplex Method of Nelder & Mead (SM) [12] and the variable metric method NEWUOA [13]. 28 independent runs of each algorithm were launched. The current optimization routine of *MatrixFrame* was not capable to yield even a rational scheme of pile placement for the problems considered in this paper.

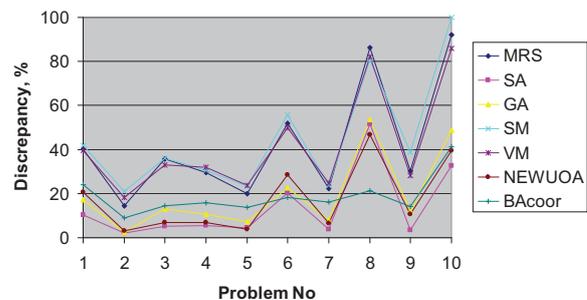


Figure 2. The best values of seven algorithms MRS, SA, GA, SM, VM, NEWUOA, and BAcoor found in 28 runs normalized to R_{ideal} (deviations in %).

In [1], only for two problems the engineering solution was achieved in 5000 objective function evaluations. However, in six problems the discrepancies between obtained and engineering solutions are 2.2 – 5.5 %.

Figure 2 illustrates the comparison of the new BAcoor algorithm with six well-known methods. The exact global solutions were not found. However, in seven examples, the deviations of obtained results from exact solutions do not exceed 5% and can be regarded as insignificant for engineering purposes. In three examples, the deviations of the best algorithms reach 20-47%. The differences of computing times were about 20-30 %.

Thus, the simulated annealing (SA) and NEWUOA outperform all the other algorithms in orderly problems; genetic algorithm (GA) is not far behind.

In this investigation ten pile placement schemes were optimized using BAcoor. To have a fair comparison of the results, objective functions were evaluated 5000 times, and 28 independent runs of the algorithm were performed. BAcoor has a clear advantage compared, for example, with GA: it has the only parameter ε to be adjusted numerically to the particular problem. Thus, for the problems No 2 and 6 it was found that $\varepsilon = 10^{-7}$, while for the remaining problems its favorable magnitude is 10^{-5} . All numerical results of 28 independent experiments are rendered in Appendix 2.

Thus, BAcoor is less efficient than SA, GA, and NEWUOA, but outperforms the remaining algorithms. However, in the awkward problems No. 8 and 6, BAcoor achieves better solutions. In the most awkward problem No. 8, it reduces the deviation to 21% (the best deviation 38% of the remaining methods was achieved by NEWUOA). A closer view to the topologies of grillages No. 8 and 6 reveals that there the piles tend to group unevenly under few connecting beams while under some beams there are

any supporting piles. Possibly, BAcoor better copes with the problem of physical discontinuity mentioned in Section 3 of the paper. Generally, in 5000 evaluations of the objective function the satisfactory solution is not achieved for all problems; more evaluations or more independent runs of an algorithm from random starting points are needed.

Summing-up of numerical results is given in Table 1 which illustrates the advantages and disadvantages of BAcoor among other explored algorithms. Appendix 3 presents graphical schemes of best pile placement schemes found with BAcoor (piles shown in red) and found in [1] (shown in blue).

5. Conclusions

In some special cases, mainly for awkward schemes where distribution of piles is very uneven BAcoor outperforms all other explored optimization algorithms. Additional advantages of the algorithm are that it converges for any function, provides higher density of observations around the global minimum, and can be adjusted numerically to a problem by controlling the single parameter. However, in orderly pile placement schemes BAcoor yields inferior results as compared with well-known algorithms such as SA, GA, and NEWUOA. Thus, BAcoor can be regarded just as a new useful member in the family of global optimization algorithms.

Appendixes

Appendix 1. Characteristics of problems.

Problem No	N_a	L	R_{allw}	R_{ideal}
1	25	172.9	325	307.47
2	18	52.9	110	104.12
3	31	84.1	105	101.85
4	31	84.9	105	101.24
5	30	63.9	100	97.51
6	37	80.1	100	97.53
7	23	129.1	300	287.35
8	34	137.9	250	236.28
9	17	97.6	250	244.71
10	55	315.61	350	349.05

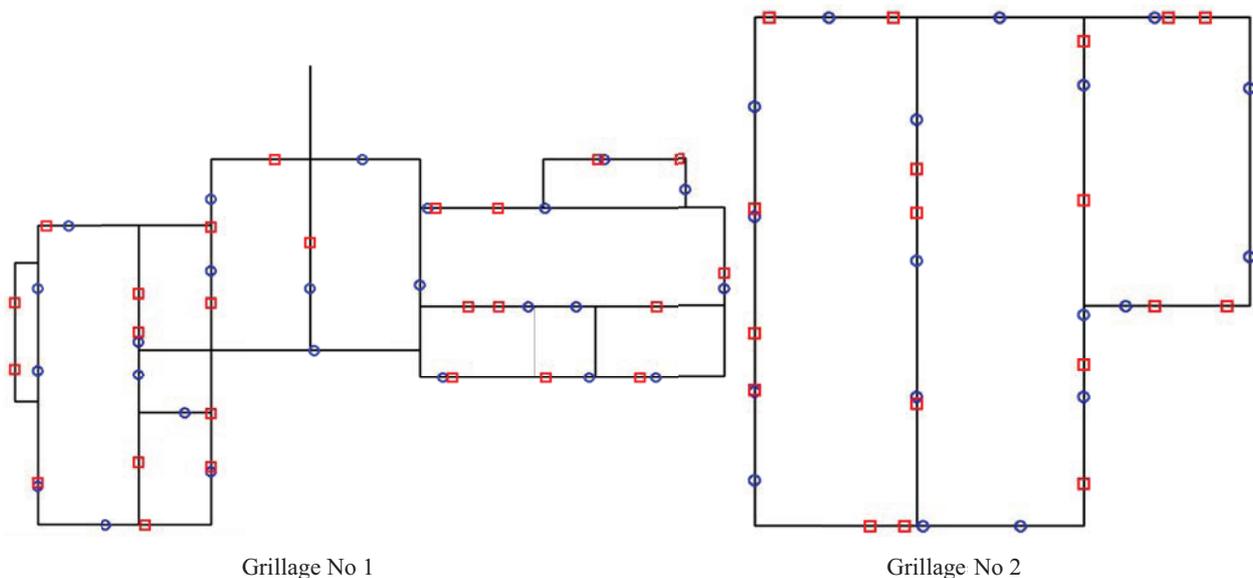
Table 1. Summary of numerical results

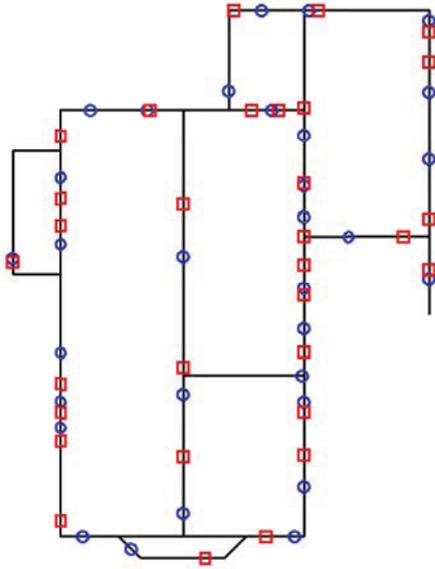
Experiment No	Best value found	Best value found in [1]	Exact solution	Place of BAcoor among algorithms
1	381.39	339.30	307.47	4
2	113.47	106.36	104.12	4
3	116.43	107.25	101.85	4
4	117.16	106.80	101.24	4
5	110.91	101.05	97.51	4
6	115.45	117.26	97.53	1
7	333.77	298.11	287.35	4
8	286.22	346.94	236.28	1
9	278.86	253.00	244.71	4
10	493.53	463.34	349.05	3

Appendix 2. Optimization results for all 10 problems in 28 independent numerical experiments (the best are highlighted in bold type)

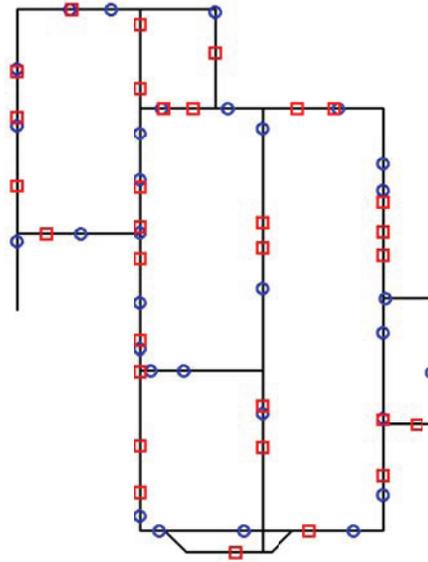
Problem / Experiment	1	2	3	4	5	6	7	8	9	10
1	422.70	123.61	133.91	124.90	118.51	151.73	370.92	495.55	310.95	652.44
2	422.32	127.64	136.68	117.99	114.60	137.42	383.00	343.94	297.16	523.63
3	381.39	118.58	118.47	122.26	119.83	129.77	362.67	516.54	316.04	548.62
4	411.51	126.89	129.21	117.97	124.48	115.45	344.09	392.20	339.15	660.61
5	419.50	121.59	124.84	124.83	122.89	146.17	362.18	464.57	301.78	531.64
6	407.47	124.19	125.42	117.27	121.33	129.75	355.68	364.42	304.55	530.02
7	385.21	122.56	127.16	119.01	122.10	130.56	364.51	330.03	301.00	523.82
8	400.56	123.45	116.43	117.16	115.78	126.63	353.76	450.64	303.93	564.01
9	404.12	128.72	125.83	121.94	114.57	129.70	370.59	382.95	317.00	634.35
10	422.51	124.55	116.57	126.10	115.00	124.10	354.68	286.22	308.70	522.27
11	431.72	125.16	121.27	119.63	116.25	132.77	341.41	362.97	335.90	634.75
12	396.67	118.55	121.20	120.63	119.65	127.03	336.15	521.20	307.87	655.23
13	416.00	130.47	123.06	125.22	114.87	124.52	360.39	353.70	306.70	493.53
14	393.81	124.77	125.30	128.50	115.37	126.34	339.22	359.17	324.96	556.71
15	388.31	113.47	121.12	120.59	123.66	130.90	352.88	412.61	308.62	525.27
16	414.56	128.69	126.91	119.71	112.90	137.12	365.99	315.30	314.54	513.90
17	423.10	122.23	132.96	127.90	117.11	128.62	352.85	336.78	340.01	659.12
18	397.76	125.79	117.31	126.44	121.86	116.41	350.32	443.44	291.91	616.27
19	510.10	127.75	121.49	123.33	113.98	129.50	357.69	511.52	321.82	586.17
20	441.61	125.65	117.72	121.16	113.64	143.30	362.40	443.02	331.86	644.46
21	401.82	121.69	124.77	120.37	118.53	132.34	345.41	306.27	326.47	625.58
22	398.30	120.43	119.73	125.72	118.81	145.54	372.29	358.19	308.28	546.57
23	398.01	116.11	123.01	120.15	120.08	130.38	352.97	344.70	313.14	601.59
24	396.72	123.88	134.05	138.72	122.44	129.42	353.85	400.11	310.54	578.04
25	390.79	120.11	128.47	119.40	122.44	133.84	341.34	422.53	303.31	542.84
26	406.33	120.66	119.04	118.86	110.91	137.25	333.77	316.98	309.55	516.35
27	398.35	124.23	124.96	179.64	113.51	130.21	348.91	337.62	278.86	613.61
28	420.52	122.32	133.19	118.23	115.88	122.55	352.10	372.81	298.18	544.10

Appendix 3. Best pile placement schemes found with BAcoor (in red color □) and found in [1] (in blue color ○).

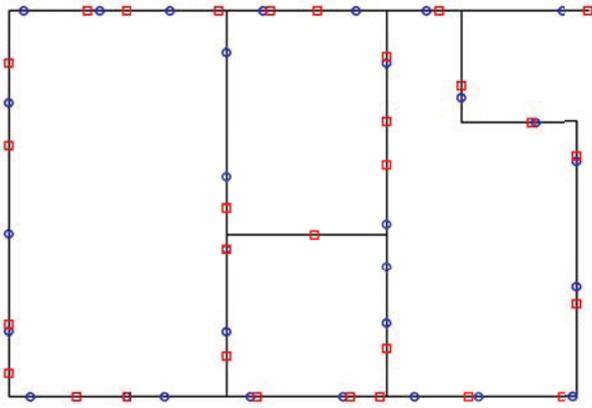




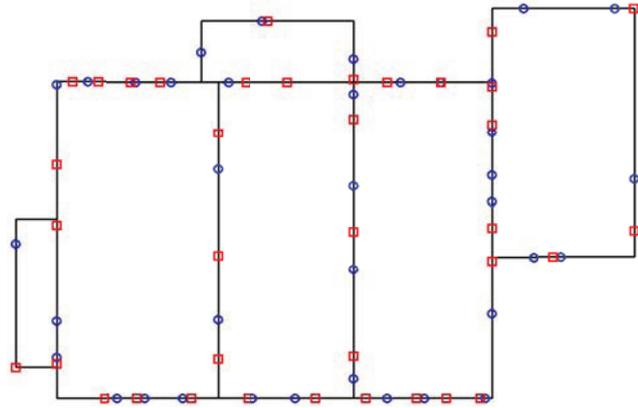
Grillage No 3



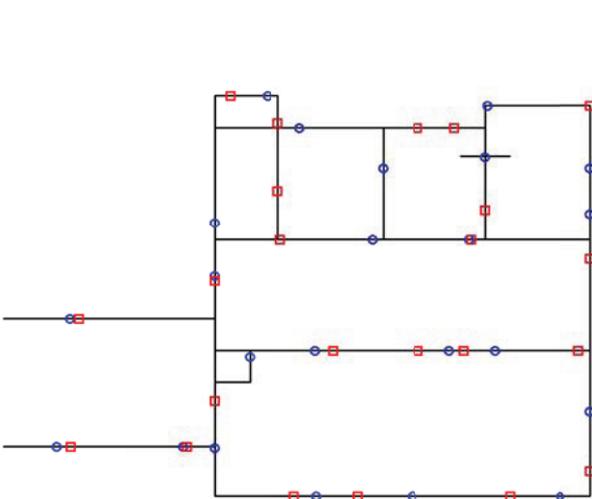
Grillage No 4



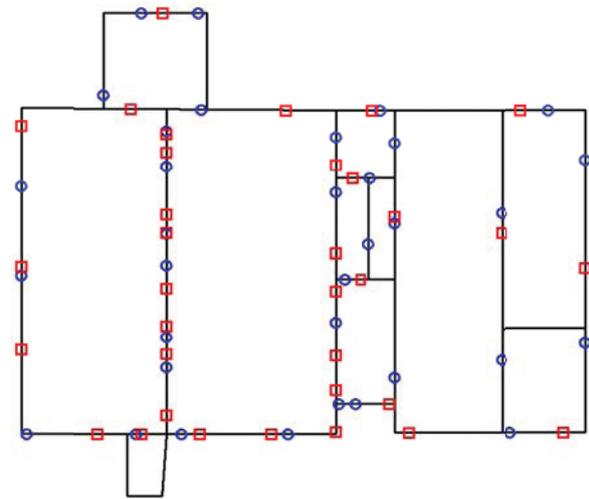
Grillage No 5



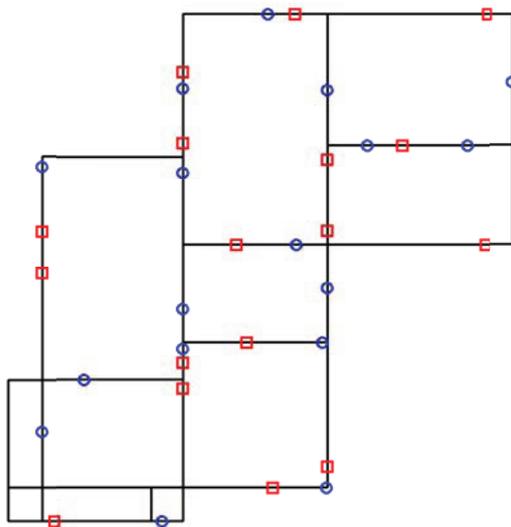
Grillage No 6



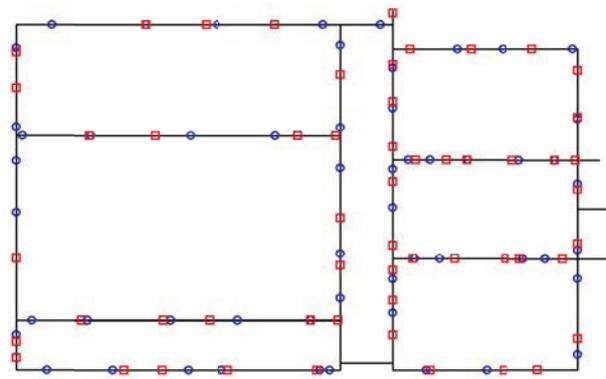
Grillage No 7



Grillage No 8



Grillage No 9



Grillage No 10

References

- [1] **R. Belevičius, S. Ivanikovas, D. Šešok, J. Žilinskas, S. Valentinavičius.** Optimal placement of piles in real grillages: experimental comparison of optimization algorithms. *Information Technology and Control*. 2011, Vol. 40, No. 2, 123 - 132.
- [2] **R. Belevičius, S. Valentinavičius.** Optimisation of grillage-type foundations. *Proceedings of 2nd European ECCOMAS and IACM Conference "Solids, Structures and Coupled Problems in Engineering", Cracow, Poland 26-29 June, 2001*, 416-421.
- [3] **R. Belevičius, S. Valentinavičius, E. Michnevich.** Multilevel optimization of grillages. *Journal of Civil Engineering and Management*. 2002, 8 (2), 98-103, <http://dx.doi.org/10.1080/13923730.2002.10531259>.
- [4] **C. M. Chan, L. M. Zhang, J. T. M. Ng.** Optimization of pile groups using hybrid genetic algorithms. *Journal of Geotechnical and Geoenvironmental Engineering*. 2009, Vol. 135, No. 4, pp. 497-505, [http://dx.doi.org/10.1061/\(ASCE\)1090-0241\(2009\)135:4\(497\)](http://dx.doi.org/10.1061/(ASCE)1090-0241(2009)135:4(497)).
- [5] **P. Diaconis.** Bayesian numerical analysis. In *Statistical Decision Theory and Related Topics*, Springer Verlag, 1988, pp. 163-175, http://dx.doi.org/10.1007/978-1-4613-8768-8_20.
- [6] **D. E. Goldberg.** Genetic Algorithms in Search, Optimization and Machine Learning, Kluwer Academic Publishers, Boston, MA, 1989, 432 p.
- [7] **K. N. Kim, S.-H. Lee, K.-S. Kim, C.-K. Chung, M. M. Kim, H. S. Lee.** Optimal pile arrangement for minimizing differential settlements in piled raft foundations. In: *Computers and Geotechnics*. 2001, 28 (4), pp. 235-253, [http://dx.doi.org/10.1016/S0266-352X\(01\)00002-7](http://dx.doi.org/10.1016/S0266-352X(01)00002-7).
- [8] **H. J. Kushner.** A new method of locating the maximum point of an arbitrary multi-peak curve in the presence of noise. *Journal of Basic Engineering*. 1964, 86: 97-100, <http://dx.doi.org/10.1115/1.3653121>.
- [9] **N. Metropolis, A. W. Rosenbluth, A. H. Teller, M. N. Rosenbluth, E. Teller.** Equation of State Calculations by Fast Computing Machines. *The Journal of Chemical Physics*. 1953, 21 (6): 1087, <http://dx.doi.org/10.1063/1.1699114>.
- [10] **J. Mockus.** Bayesian approach to global optimization. *Kluwer Academic Publishers, Dordrecht-London-Boston*, 1989, 270 p., <http://dx.doi.org/10.1007/978-94-009-0909-0>
- [11] **J. Mockus, W. Eddy, A. Mockus, L. Mockus, G. Reklaitis.** Bayesian Heuristic Approach to Discrete and Global Optimization. *Kluwer Academic Publishers, ISBN 0-7923-4327-1, Dordrecht-London-Boston*, 1997, 416 p.
- [12] **J. A. Nelder, R. Mead.** (1965). A simplex method for function minimization. *Computer Journal*. 1965, 7: 308-313.
- [13] **M. J. D. Powell,** The NEWUOA software for unconstrained optimization without derivatives. In: *Di Pillo, G., Roma, M. (eds) Large-Scale Nonlinear Optimization. Vol. 83 of Nonconvex Optimization and Its Applications*, Springer, 2006, 255-296.
- [14] **L. C. Reese, W. M. Isenhower, S.-T. Wang.** Analysis and Design of Shallow and Deep Foundations. *John Wiley & Sons*, 2006, pp. 608.
- [15] **V. R. Šaltenis.** On a method of multi-extremal optimization. *Automatics and Computers (Avtomatika i Vychislitel'naya Tekhnika)*. 1971, (3): 33-38. (in Russian).
- [16] **A. Torn, A. Žilinskas.** Global optimization. Springer-Verlag, Berlin, 1989, 255, <http://dx.doi.org/10.1007/3-540-50871-6>.

Received May 2012.